



Technische
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Introduction to IT-Security WS 20/21

Exercise 01

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1 Security goals

2 Simple combinatorics

3 XOR cipher

Following the rule $a \oplus b = c \Leftrightarrow b = a \oplus c$ these calculations can be made:

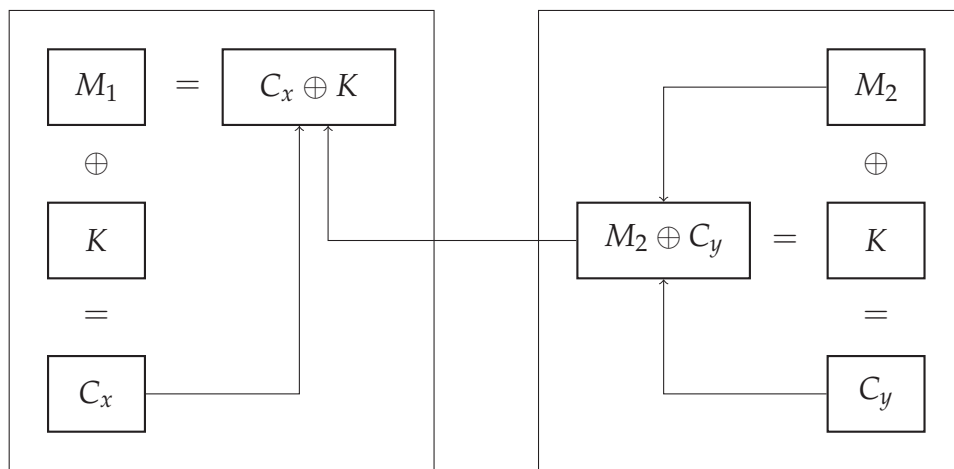


Figure 1: This Diagram shows how an attacker can calculate the key K and the message M_1 . C_x , C_y and M_2 are known to the attacker.

A few requirements must be satisfied in order to get hold of K and M_1 :

- M_2 must be longer than M_1 or K , so that the key can be calculated in at least the needed length.
- A successfully decoded message must be distinguishable from an unsuccessfully decoded message, so that the cipher texts C_x and C_y can be exchanged if necessary.